

Dependency relations, modalities and the syntax-semantics interface

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Abstract

The type logics $(N)L(P)_{\diamond}$ extend Lambek's categorial systems with adjoint pairs of unary modalities \Diamond, \Box . In previous work, modalities have been used as licensors for controlled forms of restructuring, reordering and copying.

Here, we study a complementary use of the modalities as dependency features coding for grammatical roles. The result is a multidimensional type logic simultaneously inducing dependency and function argument structure on the linguistic material.

Background: NWO project "A composition calculus for vector-based semantic modelling with a localization for Dutch". Thanks to Konstantinos Kogkalidis, Gijs Wijnholds.

Outline

- ▶ Recap: modalities for structural control
- ▶ A dependency-enhanced type logic: theory and practice
- ▶ Challenges: integrating logical and structural reasoning

Legend Core of Lambek CG: residuated families of type-forming operations.

$$\begin{array}{c} A \rightarrow C/B \quad \text{iff} \quad A \bullet B \rightarrow C \quad \text{iff} \quad B \rightarrow A \setminus C \\ \Diamond A \rightarrow B \quad \text{iff} \quad A \rightarrow \Box B \end{array}$$

Extensions: global, or \Diamond -controlled, structural options

- **NL**, pure residuation logic Lambek 1961
- **L**, the former plus global associativity Lambek 1958
- **LP**, aka MILL, associative+commutative Lambek-Van Benthem
- subexponentials for restricted copying, McPheat, Valentin, Kuznetsov @MALIN

Note on proof format, compositionality

N.D., sequent-style Judgements $X \vdash A$, formula A , structure X : $X, Y ::= A \mid X \cdot Y$.

Notation $X[Y]$ substructure Y in context X . Axioms $A \vdash A$

Logical rules, e.g. \ Elimination, Introduction

$$\frac{X \vdash A \quad Y \vdash A \setminus B}{X \cdot Y \vdash B} \setminus E \quad \frac{A \cdot X \vdash B}{X \vdash A \setminus B} \setminus I$$

Explicit structural rules, e.g. right / left rotation of structure tree

$$\frac{X[Y \cdot (Z \cdot W)] \vdash D}{X[(Y \cdot Z) \cdot W] \vdash D} A^r \quad \frac{X[(Y \cdot Z) \cdot W] \vdash D}{X[Y \cdot (Z \cdot W)] \vdash D} A^l$$

Steno words instead of their types left of turnstile; Ax: lex type assignment.

$$\frac{\begin{array}{c} \text{the} \\ \hline np/n \end{array} \quad \begin{array}{c} \text{temperature} \\ \hline n \end{array}}{\text{the} \cdot \text{temperature} \vdash np} /E \quad \frac{\begin{array}{c} \text{rises} \\ \hline np \setminus s \end{array}}{(the \cdot \text{temperature}) \cdot \text{rises} \vdash s} \setminus E$$

Compositionality

Montague's view Compositional interpretation as a structure-preserving map

[·] : Source → Target ‘Universal Grammar’, 1970

sending types and proofs of a source logic to their target counterparts.

Concretely Source (N)LP+, Target MILL/LP

$$X_{[A_1, \dots, A_n]} \vdash_{(\mathbf{N})\mathbf{L}_+} B \quad \xrightarrow{\lceil \cdot \rceil} \quad x_1 : \lceil A_1 \rceil, \dots, x_n : \lceil A_n \rceil \vdash_{\text{MILL}/\mathbf{LP}} M : \lceil B \rceil$$

Shortcut LP/MILL target terms associated with proofs in source logic (**N**)**L+**.

$$\frac{X \vdash N^{[A]} : A \quad Y \vdash M^{[A] \multimap [B]} : A \setminus B}{X \cdot Y \vdash M^{[A] \multimap [B]} \ N^{[A]} : B} E \setminus \quad \frac{x^{[A]} : A \cdot X \vdash M^{[B]} : B}{X \vdash \lambda x^{[A]}.M^{[B]} : A \setminus B} I \setminus$$

Alternative target sCCC, symmetric Compact Closed Category, concretely **FdVect**. Direct, or via linear λ Coecke et al 2013, Wijnholds & Sadrzadeh 2019

Recall: the need for control

(N)L_◇: the argument

- ▶ languages exhibit phenomena that seem to require some form of
reordering, restructuring, copying
- ▶ *global* structural options are problematic
too little (undergeneration), too much (overgeneration)
- ▶ extended type language with operations for structural control:
 - ▷ *licensing* structural reasoning that is lacking by default
 - ▷ *blocking* structural reasoning that would otherwise be available

Associativity: too little

$$\frac{\text{what}}{\underline{np/(s/np)}} \quad \frac{\text{Maisie}}{np} \quad \frac{\text{knew}}{(np \setminus s)/np} \quad np \vdash np^1$$
$$\frac{}{\frac{}{\frac{\text{knew} \cdot np \vdash np \setminus s}{\text{Maisie} \cdot (\text{knew} \cdot np) \vdash s} /E} \backslash E} A^r$$
$$\frac{\text{Maisie} \cdot (\text{knew} \cdot np) \vdash s}{(\text{Maisie} \cdot \text{knew}) \cdot np \vdash s} /I^1$$
$$\frac{}{\frac{\text{Maisie} \cdot \text{knew} \vdash s/np}{\text{what} \cdot (\text{Maisie} \cdot \text{knew}) \vdash np} /E} /E$$

Compare

► what Maisie knew ..

position of the hypothesis reachable thanks to A^r (right rotation)

☺

► what Maisie knew .. about her parents

too little: A^r doesn't give access to an internal gap

☺

Associativity: too much

Locality constraint himself :: $((np \setminus s)/np) \setminus (np^\sigma \setminus s)$, herself :: $((np \setminus s)/np) \setminus (np^\Omega \setminus s)$

Alice hurt herself / *himself

Alice thinks Bob hurt *herself / himself

No way of distinguishing simple transitive verb $(np \setminus s)/np \neq$ string reducing to $(np \setminus s)/np$

$$\frac{\text{Alice} \quad np^\Omega}{\frac{\text{thinks} \quad \frac{\text{Bob} \quad \frac{\text{hurt}}{(np \setminus s)/np \quad np \vdash np^1} /E}{\frac{\text{hurt} \cdot np \vdash np \setminus s}{\frac{\text{Bob} \cdot (\text{hurt} \cdot np) \vdash s}{\frac{\text{thinks} \cdot (\text{Bob} \cdot (\text{hurt} \cdot np)) \vdash np \setminus s}{\frac{\text{thinks} \cdot ((\text{Bob} \cdot \text{hurt}) \cdot np) \vdash np \setminus s}{\frac{\text{thinks} \cdot (\text{Bob} \cdot \text{hurt}) \cdot np \vdash np \setminus s}{\frac{\text{thinks} \cdot (\text{Bob} \cdot \text{hurt}) \vdash (np \setminus s)/np}{\frac{}{(thinks \cdot (\text{Bob} \cdot \text{hurt})) \cdot \text{herself} \vdash np^\Omega \setminus s}} /I^1 \quad \frac{\text{herself}}{((np \setminus s)/np) \setminus (np^\Omega \setminus s)}} /E}} /E}} /E}} /E$$

Modalities for structural control

- ▶ The type language is extended with a pair of unary connectives: \diamond, \Box satisfying

$$\frac{\diamond A \longrightarrow B}{A \longrightarrow \Box B}$$

- ▶ Logic: \diamond, \Box form a residuated pair. One easily shows

compositions: $\diamond\Box A \longrightarrow A$ (interior) $A \longrightarrow \Box\diamond A$ (closure)

monotonicity: from $A \longrightarrow B$ infer $\diamond A \longrightarrow \diamond B$, $\Box A \longrightarrow \Box B$

- ▶ Structure: *global* rules $\leadsto \diamond$ controlled *restricted* versions, e.g.

$$A^\diamond : (A \bullet B) \bullet \diamond C \longrightarrow A \bullet (B \bullet \diamond C)$$

$$C^\diamond : (A \bullet B) \bullet \diamond C \longrightarrow (A \bullet \diamond C) \bullet B$$

Multimodal generalization families $\{\diamond_i, \Box_i\}_{i \in I}$ for particular structural choices

Control operators: N.D. rules

Structures $X, Y ::= A \mid \langle X \rangle \mid X \cdot Y$

$$\frac{\langle X \rangle \vdash A}{X \vdash \Box A} \Box I \quad \frac{X \vdash \Box A}{\langle X \rangle \vdash A} \Box E$$
$$\frac{X \vdash A}{\langle X \rangle \vdash \Diamond A} \Diamond I \quad \frac{Y \vdash \Diamond A \quad X[\langle A \rangle] \vdash B}{X[Y] \vdash B} \Diamond E$$

Shorthand $(\Diamond E')$ if left premise of $(\Diamond E)$ is an axiom:

$$\frac{X[\langle A \rangle] \vdash B}{X[\Diamond A] \vdash B} \Diamond E'$$

Control operators: terms

Terms: $M, N ::= x \mid \lambda x. M \mid MN \mid {}^{\cup}M \mid {}^{\cap}M \mid {}^{\vee}M \mid {}^{\wedge}M$

$$\frac{\langle X \rangle \vdash M : A}{X \vdash {}^{\wedge}M : \square A} \quad \square I \qquad \frac{X \vdash M : \square A}{\langle X \rangle \vdash {}^{\vee}M : A} \quad \square E$$
$$\frac{X \vdash M : A}{\langle X \rangle \vdash {}^{\cap}M : \diamond A} \quad \diamond I \qquad \frac{Y \vdash M : \diamond A \quad X[\langle x : A \rangle] \vdash N : B}{X[Y] \vdash N[{}^{\cup}M/x] : B} \quad \diamond E$$

Proof normalization: ${}^{\vee\wedge}M = M$, ${}^{\wedge\vee}M = M$; ${}^{\cup\cap}M = M$, ${}^{\cap\cup}M = M$

Concrete interpretation See Correia et al, 2020, Putting a Spin on Language: A Quantum Interpretation of Unary Connectives for Linguistic Applications.

Controlled extraction: too little \rightsquigarrow just fine

$\Diamond \Box np$: ‘moveable’ np ; key-and-lock: contract $\Diamond \Box np$ to np , once in place.

	$\frac{\text{found} \quad \square np \vdash \square np}{\langle \square np \rangle \vdash np} \square E$
	$\frac{(np \setminus s)/np \quad \langle \square np \rangle \vdash np}{\text{found} \cdot \langle \square np \rangle \vdash np \setminus s} /E$
Alice	$\frac{\text{found} \cdot \langle \square np \rangle \vdash np \setminus s \quad (np \setminus s) \setminus (np \setminus s)}{(\text{found} \cdot \langle \square np \rangle) \cdot \text{there} \vdash np \setminus s} \setminus E$
<i>np</i>	$\frac{\text{Alice} \cdot ((\text{found} \cdot \langle \square np \rangle) \cdot \text{there}) \vdash s}{\text{Alice} \cdot ((\text{found} \cdot \text{there}) \cdot \langle \square np \rangle) \vdash s} C_{\diamond}^r$
	$\frac{\text{Alice} \cdot ((\text{found} \cdot \text{there}) \cdot \langle \square np \rangle) \vdash s}{(\text{Alice} \cdot (\text{found} \cdot \text{there})) \cdot \langle \square np \rangle \vdash s} A_{\diamond}^r$
	$\frac{(\text{Alice} \cdot (\text{found} \cdot \text{there})) \cdot \langle \square np \rangle \vdash s}{(\text{Alice} \cdot (\text{found} \cdot \text{there})) \cdot \Diamond \square np \vdash s} \Diamond E'$
what	$\frac{(\text{Alice} \cdot (\text{found} \cdot \text{there})) \cdot \Diamond \square np \vdash s}{\text{Alice} \cdot (\text{found} \cdot \text{there}) \vdash s / \Diamond \square np} /I$
<i>np/(s/◊□np)</i>	$\frac{\text{Alice} \cdot (\text{found} \cdot \text{there}) \vdash s / \Diamond \square np}{\text{what} \cdot (\text{Alice} \cdot (\text{found} \cdot \text{there})) \vdash np} /E$

$\text{A}^r_{\diamondsuit}$: controlled Associativity, $(A \bullet B) \bullet \diamondsuit C \longrightarrow A \bullet (B \bullet \diamondsuit C)$

C_\diamond^r : controlled Commutativity, $(A \bullet B) \bullet \diamond C \rightarrow (A \bullet \diamond C) \bullet B$

Island constraints: too much \rightsquigarrow just fine

Compare the following with **in**, **during** :: $(iv \setminus iv)/np$, and **which** :: $(n \setminus n)/(s/\Diamond \Box np)$

Napoleon slept in this bed
N slept during the speech
the bed which N slept in
*the speech which N slept during

In general, English allows preposition stranding, which is derivable with the controlled asso/commu rules. But some modifiers behave as *islands*, inaccessible for extraction.

\diamond **as an obstacle** a modified type assignment imposes the desired island constraint:
 $iv := np \setminus s$ **during** :: $(\Box(iv \setminus iv))/np$ Morrill 1992

- ▶ **during** first has to compose with its *np* object
- ▶ the result type $\Box(iv \setminus iv)$ is *locked* by \Box
- ▶ \Diamond *unlocks* $\Box(iv \setminus iv)$, thus sealing off **during np** as an island

We refine the idea of modalities projecting locality domains to take into account dependency info.

Dependency structure

Dependency roles articulate the linguistic material on the basis of two oppositions:

- ▶ head - **complement** relations
 - ▷ verbal domain: subj, (in)direct object, ...
 - ▷ nominal domain: prepositional object, ...
- ▶ **adjunct** - head relations
 - ▷ verbal domain: (time, manner, ...) adverbial
 - ▷ nominal domain: adjectival, numeral, determiner, ...

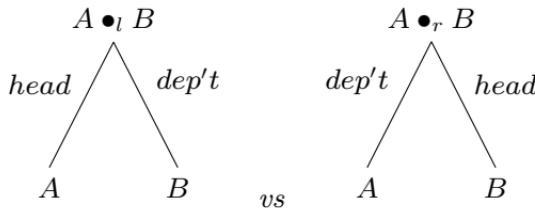
Compare: fa-structure: function vs argument

Orthogonality The fa and the dependency articulation are in general not aligned. This asks for a multidimensional type logic.

E.g. Determiner. Semantically, characteristic function of ($\llbracket N \rrbracket$, $\llbracket VP \rrbracket$) relation; morphologically, dependent on head noun.

A first step: bimodal syntactic calculus

Product \bullet is split in a left-headed \bullet_l and a right-headed \bullet_r version:



RES $A \rightarrow C /_l B \quad \text{iff} \quad A \bullet_l B \rightarrow C \quad \text{iff} \quad B \rightarrow A \setminus_l C$
 $A \rightarrow C /_r B \quad \text{iff} \quad A \bullet_r B \rightarrow C \quad \text{iff} \quad B \rightarrow A \setminus_r C$

head functor: $C /_l B, A \setminus_r C$

dependent functor: $C /_r B, A \setminus_l C$

Ref MM and Morrill 1991, Heads and phrases. Type calculus for dependency and constituent structure.

Deconstructing the headed product

Define \bullet_l, \bullet_r as compositions of regular \bullet and modal marking of the dependent:

$$A \bullet_l B := A \bullet \diamond B \quad A \bullet_r B := \diamond A \bullet B$$

Residuation: translation of the slashes

recall: $\diamond A \rightarrow B$ iff $A \rightarrow \square B$

$$\begin{array}{c} A \rightarrow C/\diamond B \\ \hline \frac{A \rightarrow C/lB}{A \bullet_l B \rightarrow C} \end{array} \quad \begin{array}{c} \frac{A \bullet \diamond B \rightarrow C}{\diamond B \rightarrow A \setminus C} \\ \hline \frac{B \rightarrow A \setminus lC}{B \rightarrow \square(A \setminus C)} \end{array} \quad \begin{array}{c} A \rightarrow C/rB \\ \hline \frac{A \bullet_r B \rightarrow C}{B \rightarrow A \setminus rC} \end{array} \quad \begin{array}{c} A \rightarrow \square(C/B) \\ \hline \frac{\diamond A \rightarrow C/B}{\diamond A \bullet B \rightarrow C} \end{array}$$

Multimodal generalization families $\{\diamond_d, \square_d\}_{d \in \text{DepLabel}}$

- ▶ $\diamond_d A \setminus C, C/\diamond_d B$ head functor assigning dependency role d to its complement
- ▶ $\square_d(A \setminus C), \square_d(C/B)$ dependent functor projecting adjunct role d

Practice

NWO project “A composition calculus for vector-based semantic modelling with a localization for Dutch”. Resources and tools for computational study of Dutch.

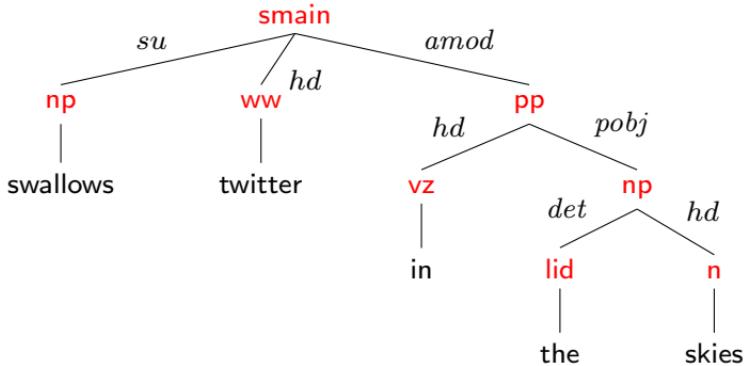
- ▶ Kogkalidis et al 2019, Constructive type-logical supertagging with self-attention networks. RepLNLP.
- ▶ Kogkalidis et al 2020a, Æthel: Automatically extracted typological derivations for Dutch. LREC.
- ▶ Kogkalidis et al 2020b, Neural proof nets. CoNLL



non-directional, dependency-enhanced types, relying on the dependency info to correctly transcribe raw text sentences into proofs and terms of the linear λ -calculus.

Structured data

Dutch treebank LASSY (google translate: zwaluwen kwetteren in 't azuur):

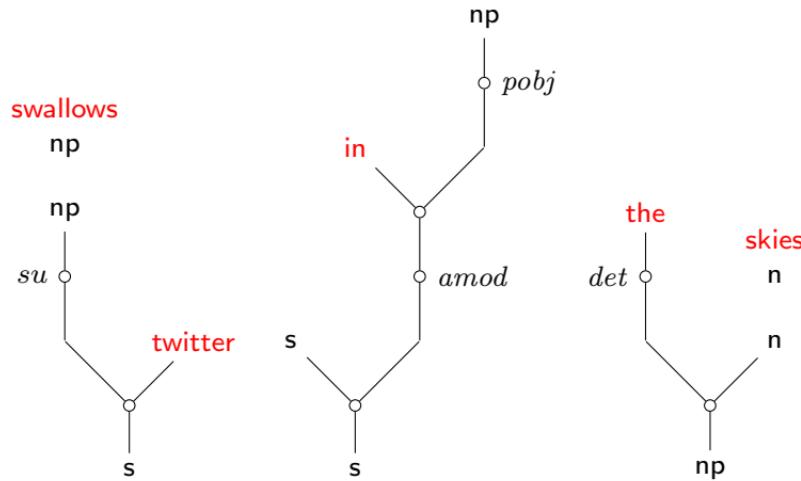


Extracted types:

swallows : np **twitter** : $\Diamond_{su} np \setminus s$ **in** : $\Box_{amod}(s \setminus s) / \Diamond_{pobj} np$ **the** : $\Box_{det}(np/n)$ **skies** : n

or non-directional: $A \setminus B, B/A \rightsquigarrow A \multimap B$

Word modules



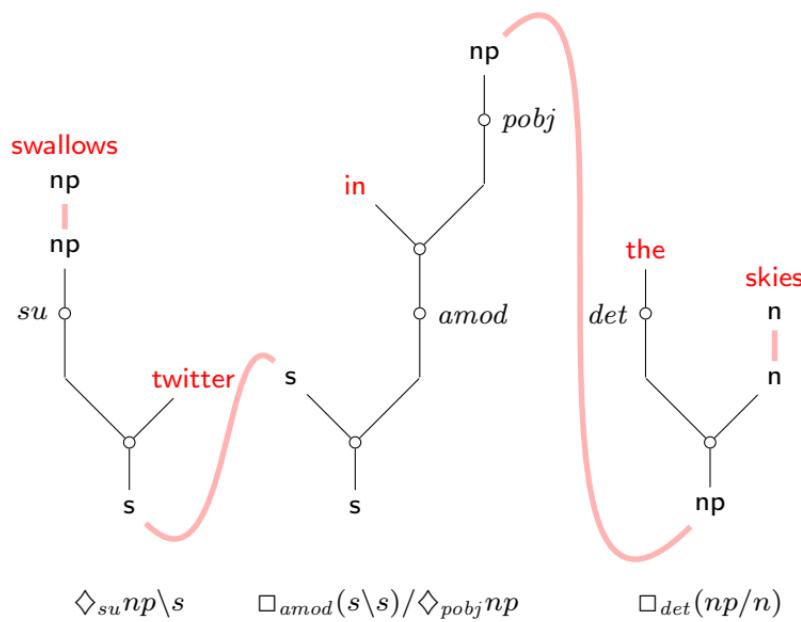
$\Diamond_{su} np \backslash s$

$\Box_{amod}(s \backslash s) / \Diamond_{pobj} np$

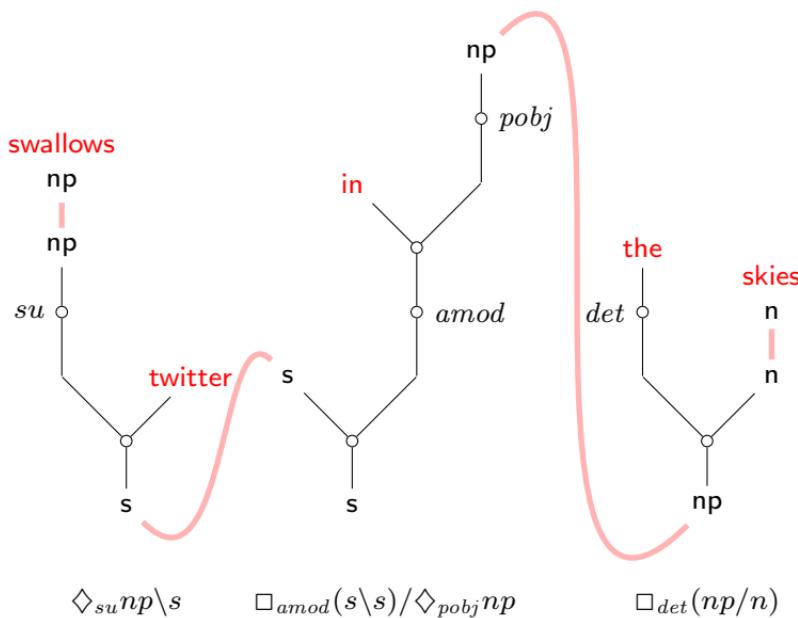
$\Box_{det}(np/n)$

Moot 2002

Word modules, proof structure, proof net



Word modules, proof net, λ term



Dependency structure

Derivation, N.D. style:

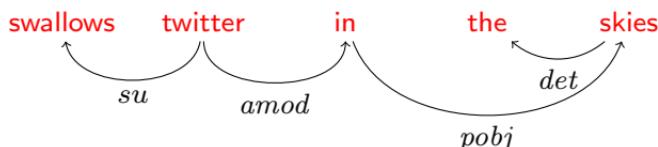
$$\begin{array}{c}
 \frac{\text{swallows}}{np} \quad \frac{\text{twitter}}{\Diamond_{sunp} \setminus s} \\
 \frac{\langle \text{swallows} \rangle^{su} \vdash \Diamond_{sunp} \quad \Diamond_{sunp} \setminus s}{\langle \text{swallows} \rangle^{su} \cdot \text{twitter} \vdash s} \quad \Diamond I \quad \backslash E \\
 \hline
 \frac{}{(\langle \text{swallows} \rangle^{su} \cdot \text{twitter}) \cdot \langle \text{in} \cdot \langle \langle \text{the} \rangle^{det} \cdot \text{skies} \rangle^{pobj} \rangle^{amod} \vdash s} \quad \Diamond I \quad \backslash E
 \end{array}$$

$$\begin{array}{c}
 \frac{\text{in}}{\Box_{amod}(s \setminus s) / \Diamond_{pobj} np} \quad \frac{\text{the}}{\Box_{det}(np/n)} \\
 \frac{}{\langle \text{the} \rangle^{det} \vdash np/n} \quad \frac{n}{\text{skies}} \\
 \hline
 \frac{\Box_{det}(np/n) \quad n}{\langle \text{the} \rangle^{det} \cdot \text{skies} \vdash np} \quad \Box E \quad / E
 \end{array}$$

$$\begin{array}{c}
 \frac{\langle \text{the} \rangle^{det} \cdot \text{skies} \vdash np}{\langle \langle \text{the} \rangle^{det} \cdot \text{skies} \rangle^{pobj} \vdash \Diamond_{pobj} np} \quad \Diamond I \\
 \hline
 \frac{\langle \langle \text{the} \rangle^{det} \cdot \text{skies} \rangle^{pobj} \vdash \Diamond_{pobj} np}{\langle \text{in} \cdot \langle \langle \text{the} \rangle^{det} \cdot \text{skies} \rangle^{pobj} \rangle^{amod} \vdash s \setminus s} \quad / E
 \end{array}$$

$$\frac{\langle \text{in} \cdot \langle \langle \text{the} \rangle^{det} \cdot \text{skies} \rangle^{pobj} \rangle^{amod} \vdash s \setminus s}{(\langle \text{swallows} \rangle^{su} \cdot \text{twitter}) \cdot \langle \text{in} \cdot \langle \langle \text{the} \rangle^{det} \cdot \text{skies} \rangle^{pobj} \rangle^{amod} \vdash s} \quad \Box E \quad \backslash E$$

Induced dependency structure:



~> within dependency domain, outgoing arcs from head to (head of) dependents

Dependency domains: some consequences

Case study: extraction revisited

Recall Dutch left-branch extraction via controlled associativity, commutativity

$$\Diamond_x A \bullet (B \bullet C) \longrightarrow (\Diamond_x A \bullet B) \bullet C \quad \Diamond_x A \bullet (B \bullet C) \longrightarrow B \bullet (\Diamond_x A \bullet C)$$

Relative pronoun **die** :: $(n \setminus n) / (!_x np \setminus s)$

$$!_x A \triangleq \Diamond_x \Box_x A$$

ambiguous between subj/obj relativization: s subordinate clause, head-final

$$\begin{array}{ll} \text{mannen die } \sqcup \text{ vrouwen haten} & \rightsquigarrow \text{men who hate women} \\ \text{mannen die vrouwen } \sqcup \text{ haten} & \rightsquigarrow \text{men who(m) women hate} \end{array}$$

Dependency refinement

- ▶ distinguish subj/obj relativisation: np hypothesis $\rightsquigarrow \Diamond_{su} np$ vs $\Diamond_{obj} np$
reduction of derivational ambiguity
- ▶ make extraction conditional on $x \succeq d$, for some boundary strength ordering
Morrill's 'island' modalities parameterized for dependency roles

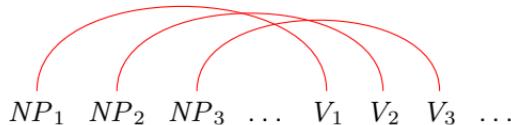
Case study: Dutch verb clusters

Crossing dependencies

Huijbregts 1984, Shieber 1985



Pattern copy language w^2 , i.e. mildly context-sensitive



Handling verb clusters: two type-logical options

- ▶ form-meaning dualism: the Abstract Categorial Grammar method

abstract syntax, divergent compositional translations:

$\lceil \cdot \rceil^{string}$ string semantics

$\lceil \cdot \rceil^{sem}$ meaning assembly

- ▶ all-in-one approach

derivations $\Gamma \vdash A$ with alternating logical and structural phases

meaning assembly: Elim/Intro logical constants

surface form: meaning preserving structural transformations

2-MCFG

- ▶ infinitival phrases $INFP$: string tuples (*rest, verb(cluster)*)
- ▶ verb raising triggers IVR_i : head-adjunction Pollard 1984

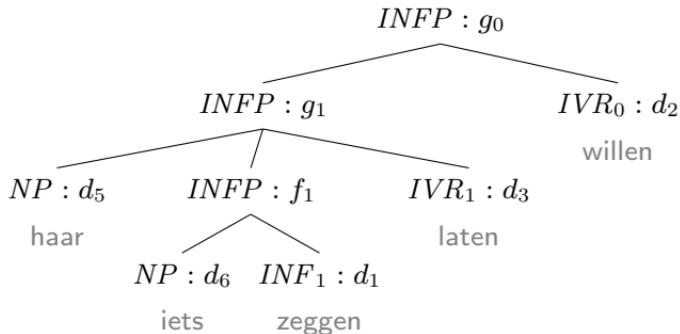
f_0	$INFP(\epsilon, x) \leftarrow INF_0(x)$	
f_1	$INFP(x, y) \leftarrow NP(x) \; INF_1(y)$	
g_0	$INFP(y, \textcolor{red}{x} \cdot z) \leftarrow INFP(y, z) \; IVR_0(\textcolor{red}{x})$	
g_1	$INFP(y \cdot z, \textcolor{red}{x} \cdot w) \leftarrow NP(y) \; INFP(z, w) \; IVR_1(\textcolor{red}{x})$	
d_0	$INF_0(\text{vertrekken}) \leftarrow$	<i>leave</i>
d_1	$INF_1(\text{zeggen}) \leftarrow$	<i>say</i>
d_2	$IVR_0(\text{willen}) \leftarrow$	<i>want</i>
d_3	$IVR_1(\text{laten}) \leftarrow$	<i>let</i>

VR triggers: modal/temporal auxiliaries (g_0); verbs of perception, causatives (g_1)

From abstract syntax to string semantics

ACG method typing rule constants, compositional translations $[\cdot] : \text{Source} \rightarrow \text{Target}$
 de Groote, Pogodalla c.s.

String semantics: higher-order modelling of tuples: $[INFP] = (\sigma \multimap \sigma \multimap \sigma) \multimap \sigma$
 $\triangleq \sigma^{(2)}$



$$\begin{aligned}
 [f_1] &= \lambda x \lambda y \lambda f. (f \ x \ y) && :: \quad \sigma \multimap \sigma \multimap \sigma^{(2)} \\
 [g_0] &= \lambda q \lambda x \lambda f. (q \ \lambda y \lambda z. (f \ y \ \textcolor{red}{x} \cdot z)) && :: \quad \sigma^{(2)} \multimap \sigma \multimap \sigma^{(2)} \\
 [g_1] &= \lambda x \lambda q \lambda y \lambda f. (q \ \lambda z \lambda w. (f \ x \cdot z \ \textcolor{red}{y} \cdot w)) && :: \quad \sigma \multimap \sigma^{(2)} \multimap \sigma \multimap \sigma^{(2)} \\
 [g_0 \ (g_1 \ d_5 \ (f_1 \ d_6 \ d_1) \ d_3) \ d_2] &= \lambda f. (f \ \text{haar} \cdot \text{iets} \ \text{willen} \cdot \text{laten} \cdot \text{zeggen})
 \end{aligned}$$

NL

The ACG method is easily adapted to a **NL** source: words as abstract constants.

	<u>iets</u>	<u>zeggen</u>		
	<u>NP</u>	<u>$NP \setminus INFP$</u>		
	$iets \cdot zeggen \vdash INFP$	$\backslash E$	<u>laten</u>	
			<u>$INFP \setminus (NP \setminus INFP)$</u>	
<u>haar</u>			$\backslash E$	
<u>NP</u>		$(iets \cdot zeggen) \cdot laten \vdash NP \setminus INFP$	$\backslash E$	
				<u>willen</u>
	$haar \cdot ((iets \cdot zeggen) \cdot laten) \vdash INFP$			<u>$INFP \setminus INFP$</u>
				$\backslash E$
	\dagger	$(haar \cdot ((iets \cdot zeggen) \cdot laten)) \cdot willen \vdash INFP$		

$\llbracket \text{zeggen} \rrbracket$	$= \lambda x \lambda f. (f\ x\ \text{zeggen})$	$:: \sigma \multimap \sigma^{(2)}$
$\llbracket \text{willen} \rrbracket$	$= \lambda q \lambda f. (q\ \lambda y \lambda z. (f\ y\ \text{willen}\cdot z))$	$:: \sigma^{(2)} \multimap \sigma^{(2)}$
$\llbracket \text{laten} \rrbracket$	$= \lambda q \lambda x \lambda f. (q\ \lambda z \lambda w. (f\ x \cdot z\ \text{laten}\cdot w))$	$:: \sigma^{(2)} \multimap \sigma \multimap \sigma^{(2)}$

compare	$\vdash^{string} = \lambda f.(f\;haar\cdot iets\;willen\cdot laten\cdot zeggen)$
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Dependency enhancement

function types $A \setminus B \rightsquigarrow \Diamond_d A \setminus B$

vc: verbal complement

$\frac{\text{iets}}{np}$	$\frac{\text{zeggen}}{\Diamond_{obj} np \setminus inf}$
	$\frac{\langle \text{iets} \rangle^{obj} \vdash \Diamond_{obj} np \quad \Diamond_{obj} np \setminus inf}{\langle \text{iets} \rangle^{obj} \cdot \text{zeggen} \vdash inf} \setminus E$
$\frac{\text{haar}}{np}$	$\frac{\langle \langle \text{iets} \rangle^{obj} \cdot \text{zeggen} \rangle^{vc} \vdash \Diamond_{vcinf}}{\Diamond_{vcinf} \setminus (\Diamond_{obj} np \setminus inf)} \Diamond I$
	$\frac{\langle \text{iets} \rangle^{obj} \cdot \text{zeggen} \vdash inf \quad \langle \langle \text{iets} \rangle^{obj} \cdot \text{zeggen} \rangle^{vc} \vdash \Diamond_{vcinf} \quad \langle \langle \text{iets} \rangle^{obj} \cdot \text{zeggen} \rangle^{vc} \cdot \text{laten} \vdash \Diamond_{obj} np \setminus inf}{\langle \langle \text{iets} \rangle^{obj} \cdot \text{zeggen} \rangle^{vc} \cdot \text{laten} \vdash \Diamond_{obj} np \setminus inf} \setminus E$
	$\frac{\langle \text{haar} \rangle^{obj} \cdot (\langle \langle \text{iets} \rangle^{obj} \cdot \text{zeggen} \rangle^{vc} \cdot \text{laten}) \vdash inf}{\langle \langle \text{haar} \rangle^{obj} \cdot (\langle \langle \text{iets} \rangle^{obj} \cdot \text{zeggen} \rangle^{vc} \cdot \text{laten}) \rangle^{vc} \vdash \Diamond_{vcinf}} \Diamond I$
	$\frac{\langle \langle \text{haar} \rangle^{obj} \cdot (\langle \langle \text{iets} \rangle^{obj} \cdot \text{zeggen} \rangle^{vc} \cdot \text{laten}) \rangle^{vc} \vdash \Diamond_{vcinf} \quad \Diamond_{vcinf} \setminus inf}{\langle \langle \text{haar} \rangle^{obj} \cdot (\langle \langle \text{iets} \rangle^{obj} \cdot \text{zeggen} \rangle^{vc} \cdot \text{laten}) \rangle^{vc} \cdot \text{willen} \vdash inf} \setminus E$

(WANT \cap ((LET \cap (SAY \cap SOMETHING)) \cap HER))

Modelling head adjunction

Logical/structural phases

$|\Gamma|$: antecedent structure term with yield Γ

⋮	<i>logical phase</i>
<hr/> $ \text{haar iets zeggen laten willen} \vdash \text{INFP}$	
<hr/> $?$	<i>structural phase</i>
<hr/> $ \text{haar iets willen laten zeggen} \vdash \text{INFP}$	

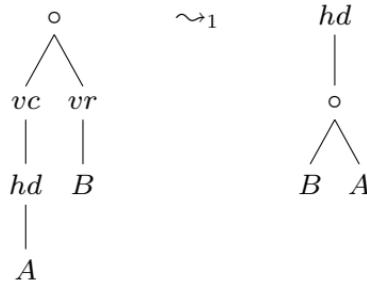
Challenge To reach the end sequent, the VR triggers have to break into the dependency domain of their infinitival complements.

As with Pollard's Head Grammars, we need **head marking**:

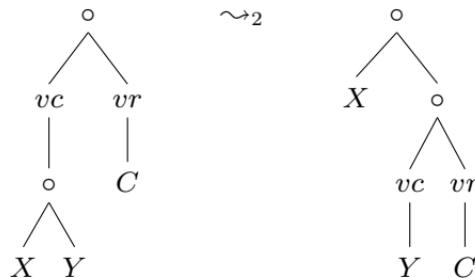
$$\begin{array}{lll} \text{zeggen} :: \Diamond_{obj} np \setminus inf & \rightsquigarrow & \Box_{hd}(\Diamond_{obj} np \setminus inf) \\ \text{willen} :: \Diamond_{vc} inf \setminus inf & \rightsquigarrow & \Box_{vr}(\Diamond_{vc} inf \setminus inf) \\ \text{laten} :: \Diamond_{vc} inf \setminus (\Diamond_{obj} np \setminus inf) & \rightsquigarrow & \Box_{vr}(\Diamond_{vc} inf \setminus (\Diamond_{obj} np \setminus inf)) \end{array}$$

Structural transformations

- ▶ *vr* trigger in construction with the head *hd* of its *vc* merges into complex *hd*:



- ▶ restructuring disassembles complex verbal complements:



iets willen zeggen

$$\frac{\text{iets}}{np} \quad \frac{\text{zeggen}}{\square_{hd}(\Diamond_{obj} np \setminus inf)} \quad \frac{}{\langle \text{zeggen} \rangle^{hd} \vdash \Diamond_{obj} np \setminus inf} \quad \square E$$

$$\frac{\Diamond I}{\langle \text{iets} \rangle^{obj} \vdash \Diamond_{obj} np} \quad \frac{}{\langle \text{iets} \rangle^{obj} \cdot \langle \text{zeggen} \rangle^{hd} \vdash inf} \quad \backslash E$$

$$\frac{\langle \text{iets} \rangle^{obj} \cdot \langle \text{zeggen} \rangle^{hd} \vdash \Diamond_{vc} inf}{\langle \langle \text{iets} \rangle^{obj} \cdot \langle \text{zeggen} \rangle^{hd} \rangle^{vc} \vdash \Diamond_{vc} inf} \quad \Diamond I$$

$$\frac{\text{willen}}{\square_{vr}(\Diamond_{vc} inf \setminus inf)} \quad \frac{}{\langle \text{willen} \rangle^{vr} \vdash \Diamond_{vc} inf \setminus inf} \quad \square E$$

$$\frac{}{\langle \text{willen} \rangle^{vr} \vdash \Diamond_{vc} inf \setminus inf} \quad \backslash E$$

$$\frac{\langle \langle \text{iets} \rangle^{obj} \cdot \langle \text{zeggen} \rangle^{hd} \rangle^{vc} \cdot \langle \text{willen} \rangle^{vr} \vdash inf}{\langle \text{iets} \rangle^{obj} \cdot (\langle \langle \text{zeggen} \rangle^{hd} \rangle^{vc} \cdot \langle \text{willen} \rangle^{vr}) \vdash inf} \quad R2$$

$$\frac{\langle \text{iets} \rangle^{obj} \cdot (\langle \langle \text{zeggen} \rangle^{hd} \rangle^{vc} \cdot \langle \text{willen} \rangle^{vr}) \vdash inf}{\langle \text{iets} \rangle^{obj} \cdot \langle \text{willen} \cdot \text{zeggen} \rangle^{hd} \vdash inf} \quad R1$$

(\vee WANT \cap (\vee SAY \cap SOMETHING))

- ▶ joint effect: verb raising trigger left-adjoined to the head of its verbal complement
- ▶ more complex clusters: $R1 \circ R2^+$

Recombinant behaviour

$$\frac{\vdots}{\langle \langle \text{haar} \rangle^{obj} \cdot (\langle \langle \langle \text{iets} \rangle^{obj} \cdot \langle \text{zeggen} \rangle^{hd} \rangle^{vc} \cdot \langle \text{laten} \rangle^{vr}) \rangle^{vc} \cdot \langle \text{willen} \rangle^{vr} \vdash \text{inf}} \text{logical rules}$$
$$\frac{\langle \text{haar} \rangle^{obj} \cdot (\langle \langle \langle \text{iets} \rangle^{obj} \cdot \langle \text{zeggen} \rangle^{hd} \rangle^{vc} \cdot \langle \text{laten} \rangle^{vr}) \rangle^{vc} \cdot \langle \text{willen} \rangle^{vr} \vdash \text{inf}}{R2}$$
$$\frac{\langle \text{haar} \rangle^{obj} \cdot (\langle \langle \langle \text{iets} \rangle^{obj} \cdot (\langle \langle \langle \text{zeggen} \rangle^{hd} \rangle^{vc} \cdot \langle \text{laten} \rangle^{vr}) \rangle^{vc} \cdot \langle \text{willen} \rangle^{vr} \vdash \text{inf}}{R2}$$
$$\frac{\langle \text{haar} \rangle^{obj} \cdot (\langle \text{iets} \rangle^{obj} \cdot (\langle \langle \langle \text{zeggen} \rangle^{hd} \rangle^{vc} \cdot \langle \text{laten} \rangle^{vr}) \rangle^{vc} \cdot \langle \text{willen} \rangle^{vr}) \vdash \text{inf}}{R2}$$
$$\frac{\langle \text{haar} \rangle^{obj} \cdot (\langle \text{iets} \rangle^{obj} \cdot (\langle \langle \langle \text{laten} \cdot \text{zeggen} \rangle^{hd} \rangle^{vc} \cdot \langle \text{willen} \rangle^{vr}) \vdash \text{inf}}{R1}$$
$$\frac{\langle \text{haar} \rangle^{obj} \cdot (\langle \text{iets} \rangle^{obj} \cdot \langle \text{willen} \cdot (\text{laten} \cdot \text{zeggen}) \rangle^{hd}) \vdash \text{inf}}{R1}$$

$(\vee \text{WANT} \cap ((\vee \text{LET} \cap (\vee \text{SAY} \cap \text{SOMETHING})) \cap \text{HER}))$

Conclusion

- ▶ modalities as licensors of structural transformations
- ▶ dependency-enhanced types: modalities demarcating domains of locality
- ▶ options for integrating logical and structural reasoning

