A combinatorial generalization of the degree of lattice polytopes

> Benjamin Nill (Case Western Reserve University) joint work with Arnau Padrol

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I. The degree of lattice polytopes - Ehrhart theory

Let $P \subset \mathbb{R}^d$ be a *d*-dimensional lattice polytope (i.e., the vertex set $\mathcal{V}(P) \subset \mathbb{Z}^d$).

Main problem: Classify Ehrhart polynomials of lattice polytopes:

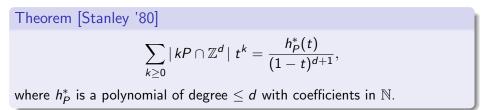
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I. The degree of lattice polytopes - Definition

Theorem [Stanley '80]

$$\sum_{k\geq 0} |\, kP\cap \mathbb{Z}^d\, |\,\, t^k = rac{h_P^*(t)}{(1-t)^{d+1}},$$

where h_P^* is a polynomial of degree $\leq d$ with coefficients in \mathbb{N} .

Definition: The degree of h_P^* is called the **degree** of *P*.

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Main problem (dimension-free version): Classify *h**-*polynomials of lattice polytopes of given degree*.

 \rightsquigarrow Degree fixed, dimension arbitrary!

I. The degree of lattice polytopes - Properties

Basic properties of the degree

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$$\deg(P) = 0 \Leftrightarrow h_P^* = 1 \Leftrightarrow P \cong \Delta_d = \operatorname{conv}(0, e_1, \dots, e_d)$$

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- $\deg(P) = 0 \Leftrightarrow h_P^* = 1 \Leftrightarrow P \cong \Delta_d = \operatorname{conv}(0, e_1, \dots, e_d)$
- deg(P) does not change by taking **lattice pyramids**.
- degree is monotone with respect to inclusion.

I. The degree of lattice polytopes - Meaning

Definition: The codegree of P:

 $\operatorname{codeg}(P) := \min\{k \in \mathbb{Z}_{>0} : \operatorname{int}(kP) \cap \mathbb{Z}^d \neq \emptyset\}$

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$$\operatorname{codeg}(P) = 1 \qquad \Leftrightarrow \quad \operatorname{deg}(P) = d \quad \Leftrightarrow \quad \operatorname{int}(P) \cap \mathbb{Z}^d \neq \emptyset$$

 $\operatorname{codeg}(P) = d + 1 \quad \Leftrightarrow \quad \operatorname{deg}(P) = 0 \quad \Leftrightarrow \quad P \cong \Delta_d$

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Definition/Lemma (Batyrev, N. '07) *P* is a **Lawrence prism**, if exists

$$\phi: \mathbb{Z}^d \twoheadrightarrow \mathbb{Z}^{d-1}, \quad \phi(P) = \Delta_{d-1}.$$

Then

$$\deg(P) \leq 1 \iff \operatorname{codeg}(P) \geq d.$$



Example: Dimension d = 2 (no interior lattice points) **Lawrence prisms**



and one exceptional triangle S

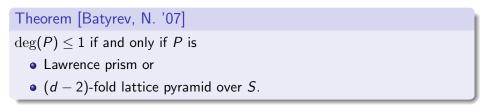


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Observe: Let $deg(P) \le 1$. Then $d \ge 3 \implies P$ Cayley polytope

I. The degree of lattice polytopes - Structure result

Theorem [Haase, N., Payne '09]

If $d > 20(\deg(P))^2$, then

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Theorem [Haase, N., Payne '09]

There exists a function f such that

$$\operatorname{Vol}_{\mathbb{Z}}(P) = 1 + h_1^* + \dots + h_{\deg(P)}^* \leq f(\deg(P), h_{\deg(P)}^*)$$

I. The degree of lattice polytopes - Cayley conjecture

Cayley conjecture If

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Refined conjecture holds, if P is smooth [Dickenstein, N. '10].

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Goal: Study (d - s)-almost neighborly polytopes for s fixed and d large!

II. The combinatorial degree of polytopes - Definition

Definition

The combinatorial codegree of d-dimensional polytope P is

 $\operatorname{codeg}_{c}(P) := \min\{|V| : V \subseteq \mathcal{V}(P), \operatorname{conv}(V) \subsetneq \partial P\}.$

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which is the maximal *codimension* of an interior face of *P*.

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If P is a lattice polytope, then

 $\deg_c(P) \leq \deg(P).$

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- Comb. degree is monotone with respect to inclusion (of vertex sets).

II. The combinatorial degree of polytopes - $\deg_c(P) = 1$

Theorem [Batyrev, N. '07]

P d-dimensional **lattice polytope** which is not a lattice pyramid. Then $deg(P) \le 1$ if and only if *P* is

- Lawrence prism or
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Theorem [N., Padrol '12]

P d-dimensional **polytope** which is not a pyramid. Then $\deg_c(P) \le 1$ if and only if *P* is

- prism over (d-1)-simplex or
- polygon.

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Definition/Proposition [N., Padrol '12]

P is a **combinatorial Cayley polytope** of r polytopes, if P is combinatorially equivalent to affine Cayley polytope.

There exists **partition** $\mathcal{V}(P) = A_1 \uplus \cdots \uplus A_r$ such that

$$\forall I \subsetneq \{1, \ldots, r\}$$
 : $\operatorname{conv}\left(\bigcup_{i \in I} A_i\right)$ is face of P .

 \Leftrightarrow

In this case, $\operatorname{codeg}_c(P) \ge r$.

Analogue to Cayley conjecture:

If $\operatorname{codeg}_c(P) > \frac{d+2}{2}$, then P is a combinatorial Cayley polytope.

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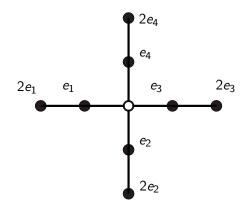
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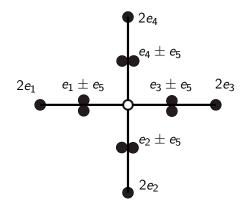
If *P* is *k*-almost neighborly with $k > \lfloor \frac{d}{2} \rfloor$, then *P* is a combinatorial Cayley polytope.

Recall:

If P is k-neighborly with $k > \lfloor \frac{d}{2} \rfloor$, then P is a simplex. II. The combinatorial degree of polytopes - Analogy? **Example:** d = 5, $\operatorname{codeg}_{c}(P) = 4 > \frac{5+2}{2}$, *P* not combinatorial Cayley



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II. The combinatorial degree of polytopes - Conjecture

Definition

P is a **weak Cayley polytope** of *r* polytopes, if there exists **cover** $\mathcal{V}(P) = A_1 \cup \cdots \cup A_r$ such that

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Correct analogue to Refined Cayley Conjecture?

If $\operatorname{codeg}_c(P) > \frac{d+r}{2}$, then P is a weak Cayley polytope of r polytopes.

 $\deg_c(P)$ is maximal codimension of interior face of *P*.

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 $\deg_c(P) = \max\{\deg(h_{\tau}) : \tau \text{ triangulation with vertex set } \mathcal{V}(P)\}$?

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Open part of Generalized Lower Bound Conjecture

P simplicial, then

 $\deg(g_P) = \min\{\deg(h_{\tau}) : \tau \text{ triangulation with vertex set } \mathcal{V}(P)\}$

where g_P is g-polynomial of face poset ∂P .