From Ehrhart theory to almost-neighborly polytopes

Benjamin Nill (Case Western Reserve University) joint work with Arnau Padrol

AMS Spring Eastern Sectional Meeting Boston College, 04/06/13 Almost-neighborly polytopes - Definition

Let P be a d-polytope.

Definition

P is k-neighborly, if any subset of k vertices is vertex set of a face.

Much studied class of polytopes (includes cyclic polytopes)

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Definition (see exercise in Grünbaum's book)

P is k-almost neighborly, if any subset of k vertices is contained in a proper face.

Not much known!

Almost-neighborly polytopes - Question Recall:

If P is k-neighborly with $k > \lfloor \frac{d}{2} \rfloor$, then P is a simplex. Almost-neighborly polytopes - Question Recall:

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Main question of this talk

What is the analogue statement for k-almost neighborly polytopes?

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What is the analogue statement for k-almost neighborly polytopes?

Example: *P* 2-almost neighborly 3-polytope, $2 > \lfloor \frac{3}{2} \rfloor$



Degree of lattice polytopes - Definition

Let $P \subset \mathbb{R}^d$ be a *d*-dimensional **lattice polytope** (i.e., the vertex set $\mathcal{V}(P) \subset \mathbb{Z}^d$).

Main object of study: Ehrhart polynomials

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Theorem [Stanley '80] $\sum_{k\geq 0} |kP \cap \mathbb{Z}^d| t^k = \frac{h_P^*(t)}{(1-t)^{d+1}},$ where h_P^* is a **polynomial of degree** $\leq \mathbf{d}$ with coefficients in \mathbb{N} .

Definition: The degree of $h_P^*(t)$ is called the **degree** of *P*.

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Problem: What can you say about lattice polytopes of given degree?

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Basic properties of the degree

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- deg(P) does not change by taking **lattice pyramids**:



• degree is monotone with respect to inclusion.

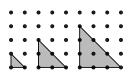
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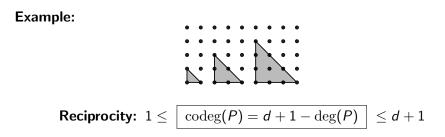
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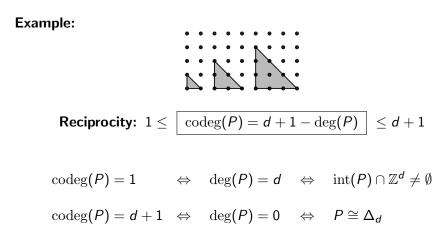
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Example: d = 2, $deg(P) \le 1$ (no interior lattice points)

Infinitely many of this type



and one exceptional triangle \boldsymbol{S}



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Examples of degree 1 in high dimensions?

Definition/Lemma (Batyrev, N. '07)

P is a Lawrence prism, if exists

$$\phi: \mathbb{Z}^d \twoheadrightarrow \mathbb{Z}^{d-1}, \quad \phi(P) = \Delta_{d-1}.$$

Then

 $\deg(P) \leq 1.$



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Theorem [Batyrev, N. '07]

P d-dimensional **lattice polytope** which is not a lattice pyramid. Then $deg(P) \le 1$ if and only if *P* is

- Lawrence prism or
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Degree of lattice polytopes - Cayley polytopes

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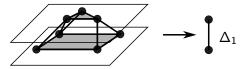
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Example: s = 2



Degree of lattice polytopes - Cayley conjecture

Corollary to Theorem:

 $\deg(P) \leq 1$ with $d > 2 \implies P$ Cayley polytope

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Cayley conjecture

 $d > 2 \deg(P) \implies P$ is Cayley polytope

Bound is **tight** (for *d* even):

 $d = 2 \operatorname{deg}(2\Delta_d)$, but $2\Delta_d$ is not Cayley polytope.

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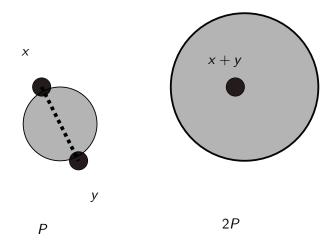
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- P is smooth \implies Cayley conjecture holds [Dickenstein, N. '10]
- $d > 20(\deg(P))^2 \Longrightarrow P$ Cayley polytope [Haase, N., Payne '09]

Combinatorial degree of polytopes - Motivation

What about the combinatorics?



Combinatorial degree of polytopes - Motivation

Observation

Let P be d-dimensional lattice polytope.

$$\operatorname{codeg}(P) > k \implies$$

P is k-almost neighborly

Definition

The combinatorial codegree of d-dimensional polytope P is

 $\operatorname{codeg}_{c}(P) := \min\{|V| : V \subseteq \mathcal{V}(P), \operatorname{conv}(V) \subsetneq \partial P\}.$

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(Co-)degree also be defined for point configurations.

Combinatorial degree of polytopes - Properties

Basic properties of the combinatorial degree

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- $\deg_c(P) = 0 \Leftrightarrow P$ is *d*-simplex
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- Comb. degree is monotone with respect to inclusion (of vertex sets).

Combinatorial degree of polytopes - $\deg_c(P) = 1$

Theorem [Batyrev, N. '07]

P d-dimensional **lattice polytope** which is not a lattice pyramid. Then $deg(P) \le 1$ if and only if *P* is

- Lawrence prism or
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Theorem [Joswig, Herrmann '10; N., Padrol '12]

P d-dimensional **polytope** which is not a pyramid. Then $\deg_c(P) \le 1$ if and only if *P* is

- prism over (d-1)-simplex or
- polygon.

One can extend this result to point configurations.

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P is a **combinatorial Cayley polytope**, if *P* is combinatorially equivalent to affine Cayley polytope

there exists **partition** $\mathcal{V}(P) = A_1 \uplus \cdots \uplus A_r$ such that

$$\forall \emptyset \neq I \subsetneq \{1, \dots, r\} : \qquad \operatorname{conv}\left(\bigcup_{i \in I} A_i\right) \text{ is proper face of } P.$$

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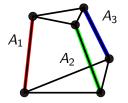
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Combinatorial Cayley conjecture?

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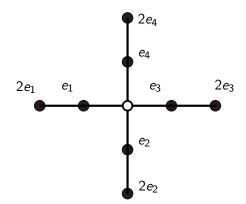
$$d > 2 \deg_c(P)$$

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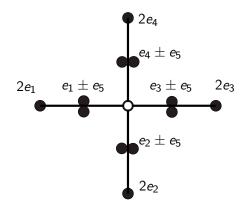
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This would solve question at beginning!

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II. The combinatorial degree of polytopes - Conjecture

Definition

P is a **weak combinatorial Cayley polytope** of *r* polytopes, if there exists *cover* $\mathcal{V}(P) = A_1 \cup \cdots \cup A_r$ such that

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Theorem [N., Padrol '12]

 $d > 3 \deg_c(P) \implies P$ is weak combinatorial Cayley polytope

Let τ be **triangulation** of *P* with vertex set $\mathcal{V}(P)$, then $\deg(h_{\tau})$ is maximal codimension of interior face of τ .

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Generalized lower bound conjecture [Murai, Nevo '12] If *P* is simplicial, then

 $\deg(g_P) = \min\{\deg(h_{\tau}) : \tau \text{ triangulation with vertex set } \mathcal{V}(P)\}$

where g_P is g-polynomial of face poset ∂P .