

Boundary Value Problems for Dirac Operators on Graphs

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Fredholm index

Definition

The (Fredholm) index of a linear operator D is given by

$$\operatorname{ind} D = \dim \ker D - \dim \ker D^*.$$

An operator is called Fredholm if this index is finite.

Index theory on closed manifolds

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- Atiyah-Singer index theorem $\implies$ index depends on the topology of the underlying manifold
- The Gauss-Bonnet theorem is a special case:

$$\text{ind}(d + d^*) = \chi(M).$$

Index theory on compact manifolds with boundary

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Consider a first-order elliptic differential operator D on a compact manifold with boundary of dimension ≥ 2 :

- D is not Fredholm if not subjected to boundary conditions
- Atiyah, Patodi and Singer found a boundary condition with respect to which D is Fredholm, and gave a formula for the index (if D is Dirac-type).

Bär-Ballmann-Bandara framework

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- Obtain a classification of all boundary conditions for which D is Fredholm.
- Simplification of the index theory for such problems.
- Led to a decomposition theorem, a general version of the relative index theorem and a generalization of the cobordism-invariance for indices.

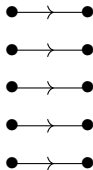
What about them?

What about one-dimensional manifolds with boundary?

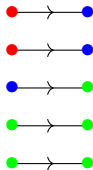
Bär-Ballmann-Bandara on 1-dim manifolds?



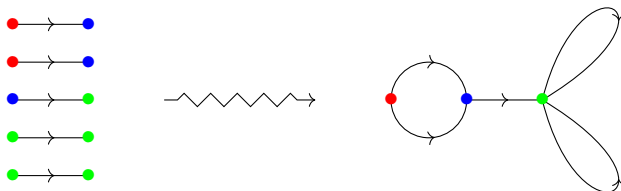
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- The spectral theory of Schrödinger operators on graphs is a classical subject with applications in physics, engineering and combinatorics.
- In order to account for the effect of spin on the quantum mechanics of graphs, first order differential operators, such as the Dirac operator, have also been studied on graphs.
- Most existing approaches focus on the spectral analysis of specific operators.

Aim/Motivation

Look at boundary value problems for metric graphs through the lens of the Bär-Ballmann framework.

In order to define first order differential operators we need to parametrize each edge, which is equivalent to

- an orientation on each edge,
- the assignment of a length to each edge.

We allow for parallel edges and loops. We thus work on finite **metric multidigraphs**.

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$$L^2(G) := \bigoplus_{e \in E} L^2([0, l_e])$$

- There is a trace map $tr : H^1(G, \mathbb{C}^r) \rightarrow \mathbb{C}^{2rn}$ mapping $f \in H^1(G, \mathbb{C}^r)$ to

$$tr(f) = (f_{e_1}(0), \dots, f_{e_n}(0), f_{e_1}(l_{e_1}), \dots, f_{e_n}(l_{e_n}))$$

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$$U_v := \left\{ (v_1, \dots, v_n, w_1, \dots, w_n) \mid \begin{array}{l} v_i = 0 \text{ if } \partial_- e_i \neq v \\ w_i = 0 \text{ if } \partial_+ e_i \neq v \end{array} \right\} \subseteq \mathbb{C}^{2rn}$$

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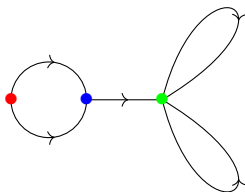
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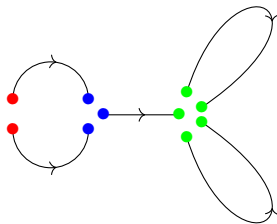
Definition

A boundary condition B is local if for each $v \in V$ there exists a subspace $B_v \subseteq U_v$ such that $B = \bigoplus_{v \in V} B_v$

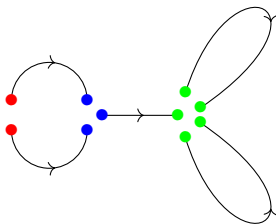
Local boundary conditions



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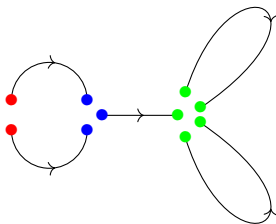


Local boundary conditions



$$B_1 \subseteq \mathbb{C}^2 \quad B_2 \subseteq \mathbb{C}^3 \quad B_3 \subseteq \mathbb{C}^5$$

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Example: The matching condition imposes that the values of a function at a vertex all agree.

Differential operators & boundary conditions

- Let $D : H^1(G, \mathbb{C}^r) \rightarrow L^2(G, \mathbb{C}^r)$ be a elliptic differential operator of first order, i.e.

$$D = \sigma \cdot d/dt + V, \text{ where } \sigma, V \in C^\infty(G, \mathbb{C}^{r \times r})$$

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- For a boundary condition B denote the restriction of D to $H_B^1(G, \mathbb{C}^r)$ by D_B
- Let $D_{\min} = D_0$ and $D_{\max} = D_{\mathbb{C}^{2m}}$.

Adjoint boundary conditions

The principal symbol σ of D defines a map $\sigma_0 : \mathbb{C}^{2rn} \rightarrow \mathbb{C}^{2rn}$

$$\begin{aligned} \sigma_0(v_1, \dots, v_n, w_1, \dots, w_n) = \\ (\sigma_{e_1}(0)v_1, \dots, \sigma_{e_n}(0)v_n, \sigma_{e_1}(l_{e_1})w_1, \dots, \sigma_{e_n}(l_{e_n})w_n). \end{aligned}$$

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Definition

For a boundary condition B define the **adjoint boundary condition**

$$B^{\text{ad}} := (\sigma_0 B)^\perp.$$

Adjoint boundary conditions II

Proposition

The adjoint of D_B is $(D^\dagger)_{B^{\text{ad}}}$, where $D^\dagger$ is the formal adjoint of D , we have $(D_{\max})^ = (D^\dagger)_{\min}$ and $(D_{\min})^* = (D^\dagger)_{\max}$.*

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- By standard ODE theory, $D_{\max}$ is surjective and $\dim \ker D_{\max} = rn$
- $\operatorname{ind} D_{\max} = rn$ and $\operatorname{ind} D_{\min} = -rn$

Index theory

Lemma

Let B be a subspace of $\mathbb{C}^{2rn}$, $B^\perp$ be the orthogonal complement to B and let $P : \mathbb{C}^{2rn} \rightarrow \mathbb{C}^{2rn}$ be the orthogonal projection onto $B^\perp$. Then,

$$L : H^1(G, \mathbb{C}^r) \rightarrow L^2(G, \mathbb{C}^r) \oplus B^\perp, \quad g \mapsto (D_{\max} g, P \operatorname{tr} g)$$

is a Fredholm operator with the same index as D_B .

The index formula

Theorem

Let $B_1 \subseteq B_2 \subseteq \mathbb{C}^{2rn}$ be boundary conditions. Then,

$$\operatorname{ind} D_{B_2} = \operatorname{ind} D_{B_1} + \dim B_2/B_1.$$

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 & L^2(G, \mathbb{C}^r) \oplus B_2^\perp & \\
 (D, P_2) \nearrow & \downarrow \text{id} \oplus \iota & \\
 H^1(G, \mathbb{C}^r) & & \\
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$$\begin{aligned}
 \text{ind } D_{B_1} &= \text{ind}(D, Q_1) = \text{ind}(D, Q_2) + \text{ind } \text{id} \oplus \iota \\
 &= \text{ind } D_{B_2} - \dim B_2 / B_1.
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Let B be any boundary condition, the index of D_B is given by

$$\operatorname{ind}(D_B) = \dim B - r|E|. \quad (1)$$

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Corollary

The index of D with respect to M is equal to the Euler characteristic of the graph, i.e.

$$\operatorname{ind} D_M = |V| - |E| = \chi(G).$$

The scalar Dirac operator and graph-like boundary conditions

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We say that A is local if $\Gamma(A)$ is a local boundary condition.

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- The existence of a invertible local map A implies that at every vertex there exist as many ingoing edges as outgoing ones.
- Local self-adjoint boundary conditions therefore exist only on Eulerian digraphs.

The spectrum of graph-like boundary conditions

- For $A : \mathbb{C}^n \rightarrow \mathbb{C}^n$, define the characteristic multinomial

$$P_A(x_1, \dots, x_n) = \det \left(\begin{pmatrix} x_1 & & \\ & \ddots & \\ & & x_n \end{pmatrix} - A \right).$$

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- The spectrum of $D_{\Gamma(A)}$ consists of all $\lambda \in \mathbb{C}$ such that

$$P_A(e^{i\lambda l_1}, \dots, e^{i\lambda l_n}) = 0,$$

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- If $l_1 = \dots = l_n = l > 0$, then the spectrum of $D_{\Gamma(A)}$ is equal to

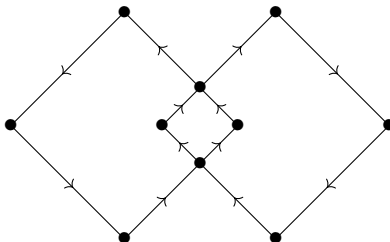
$$\bigcup_{i=1}^m \{(\arg(\lambda_i) - i \log(|\lambda_i|) + 2\pi k)/l \mid k \in \mathbb{Z}\},$$

where $\lambda_1, \dots, \lambda_m$ are the non-zero eigenvalues of A

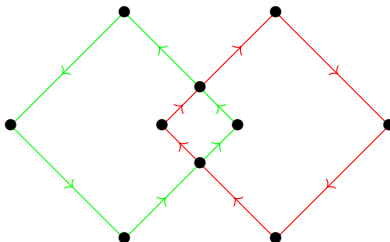
Which ones?

Which boundary conditions can we impose?

Eulerian digraphs and trail decompositions



Eulerian digraphs and trail decompositions



Examples for boundary conditions: Decompositions of Eulerian digraphs

- Every decomposition of a graph into closed directed trails yields a local permutation matrix

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- In fact, there is a 1:1 correspondence between decompositions into closed directed trails and local permutation matrices
- The spectrum of $D_{\Gamma(P)}$ for a local permutation matrix P is given by

$$\bigcup_{i=1}^m \{2\pi k/L_i \mid k \in \mathbb{Z}\},$$

where $L_1, \dots, L_m$ are the lengths of the closed trails in the corresponding decomposition.

The directed incidence matrix

Definition

The directed incidence matrix $I(G)$ of a multidigraph G is given by

$$I(G)_{ij} = \begin{cases} 1 & \text{if } \partial_+ e_j = \partial_- e_i, \\ 0 & \text{otherwise.} \end{cases}$$

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- $I(G)$ defines a local boundary condition, which in general is not self-adjoint.
- The spectrum of $D_{\Gamma(I(G))}$ is determined by the characteristic multinomial of $I(G)$, if all edges have the length 1, then it is determined by it the characteristic polynomial.

The characteristic multinomial of $I(G)$

- Let $\mathcal{C}(G) := \{C \subseteq E \mid C \text{ is a disjoint collection of cycles}\}$, for $C \in \mathcal{C}(G)$, let $\alpha(C)$ denote the amount of connected components of C .

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- For $C \in \mathcal{C}(G)$, put $x^C = \prod_{i: e_i \notin C} x_i$, we then have

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- For the characteristic polynomial $p_{I(G)}(t) = \sum_{k=0}^n a_k t^k$, we have

$$a_k = \sum_{C \in \mathcal{C}(G), |C|=n-k} (-1)^{\alpha(C)}$$

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- $a_{n-3} =$
 $-\#\{3\text{-cycles}\} + \#\{\text{disjoint pairs of loops and 2-cycles}\} -$
 $\#\{\text{disjoint triples of loops}\}$

Thanks for your attention!